

AURA: A programming language for authorization and audit

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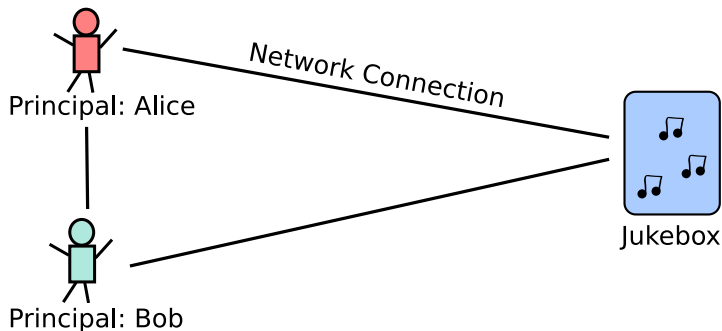
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ICFP

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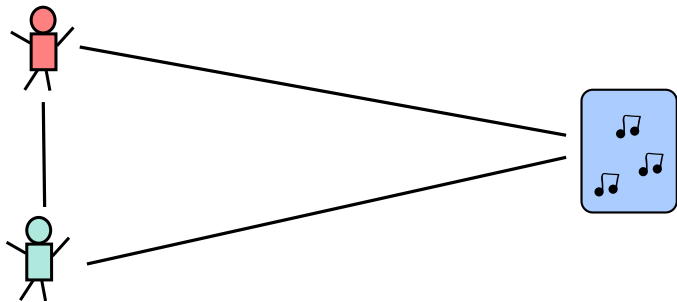
A distributed access control example



Jukebox's signature:

`playFor_raw: (s: Song) → (p: prin) → Unit`

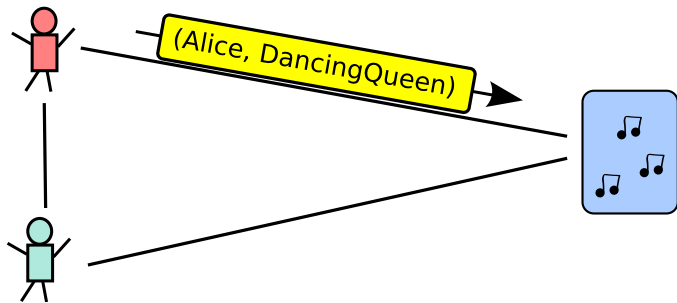
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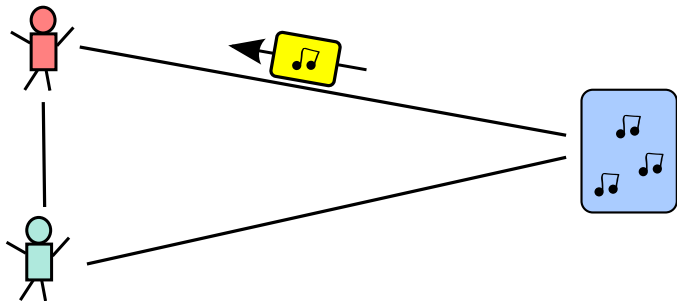
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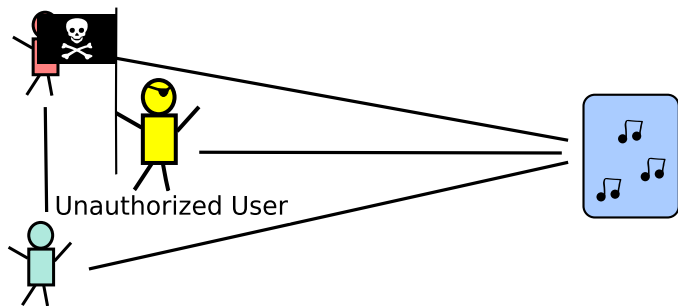
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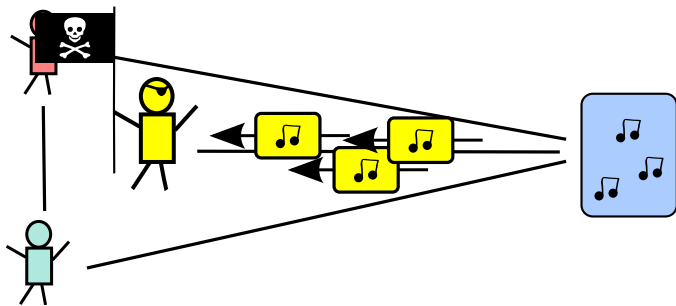
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Policy Statement (Simple):

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Policy Enforcement Problems (Hard):

- distributed decision making
- mutual distrust
- prominent use of delegation

AURA: Enforce policy with proof carrying access control.

- Programs build *proofs* attesting to their access rights.
- Proof components
 - standard rules of inference
 - *evidence* capturing principal intent (e.g. signatures)
- AURA runtime:
 - checks proof structure (well-typedness)
 - logs appropriate proofs for later audit



Proof Carrying Code [Necula+ 98], Grey Project [Bauer+ 05], Protocol Analysis [Fournet+ 07], Evidence-Based Audit [CSF 08]

Encoding policy at the ICFP server

```
shareRule  $\equiv$  self says (  
    (o: prin)  $\rightarrow$  (s: Song)  $\rightarrow$  (r: prin)  $\rightarrow$   
    (Owns o s)  $\rightarrow$   
    (o says (MayPlay r s))  $\rightarrow$   
    (MayPlay r s)))  
  
playFor: (s: Song)  $\rightarrow$  (p: prin)  $\rightarrow$   
    pf (self says (MayPlay p s))  $\rightarrow$  Unit
```

AURA features above: **pf**, **self**, **says**, dependency, effects. . .

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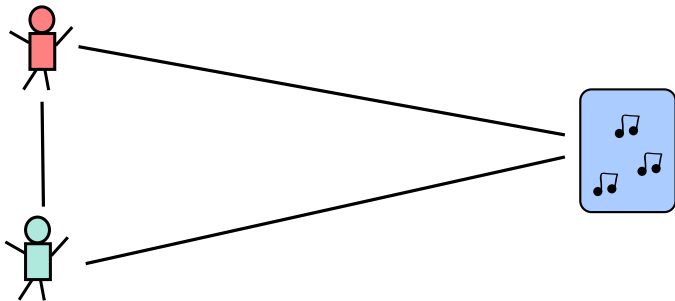
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Key Property

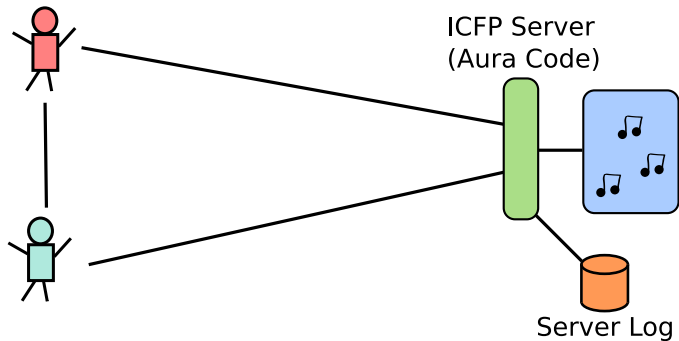
A program can only call playFor when it has an appropriate access control proof.

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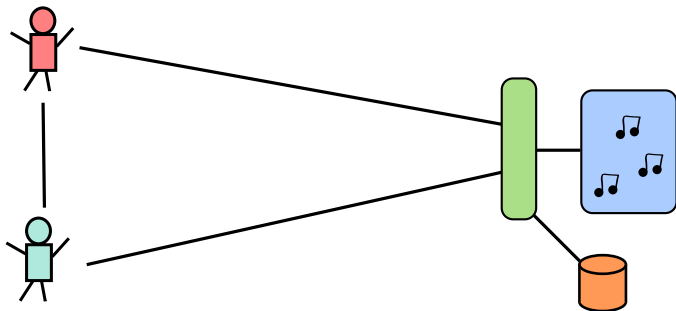
Using the ICFP policy.



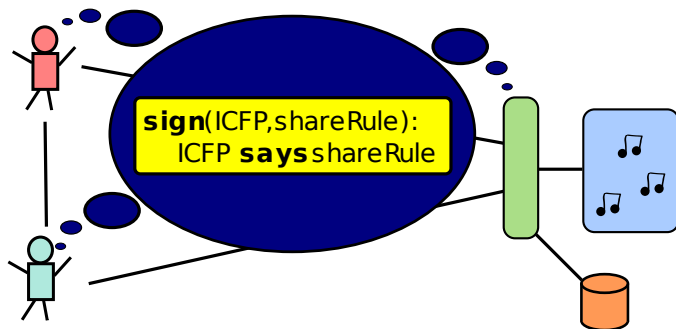
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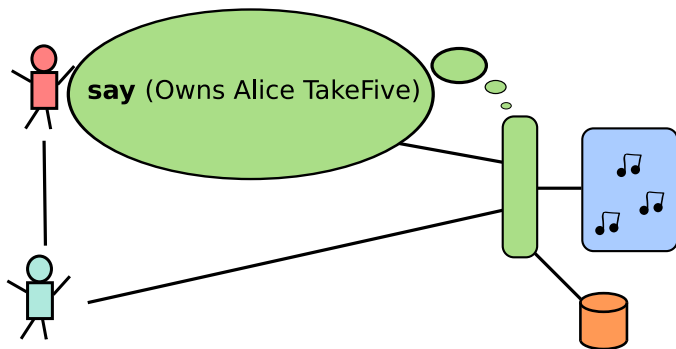
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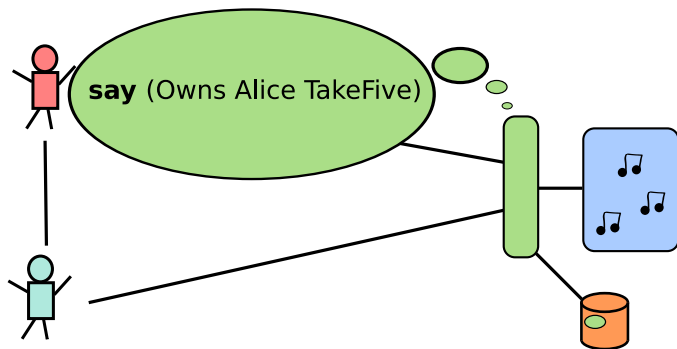
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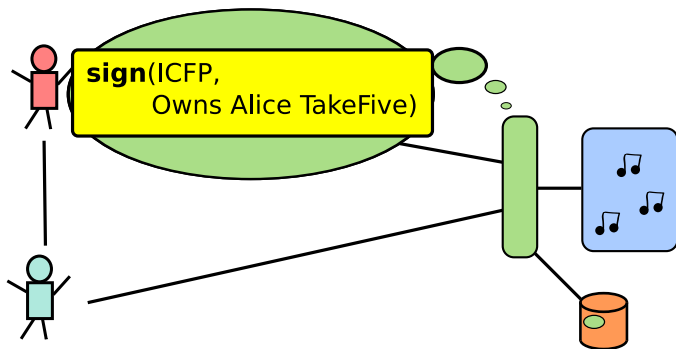
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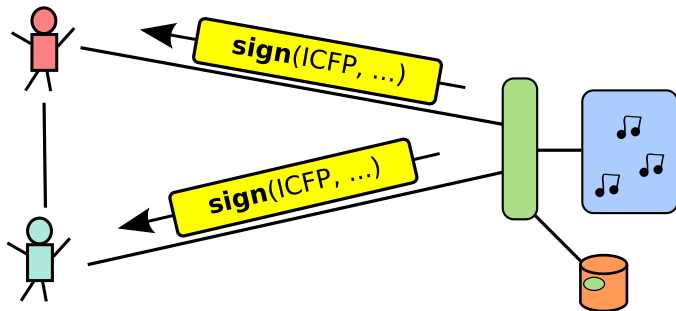
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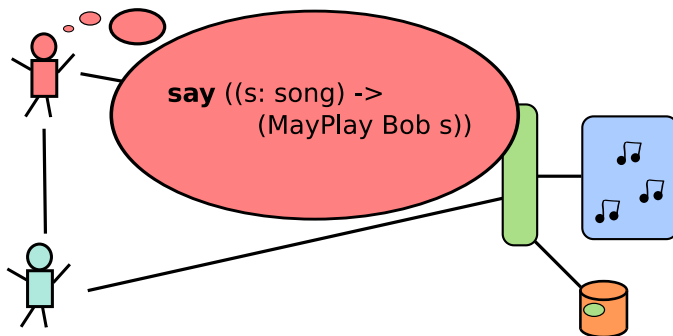
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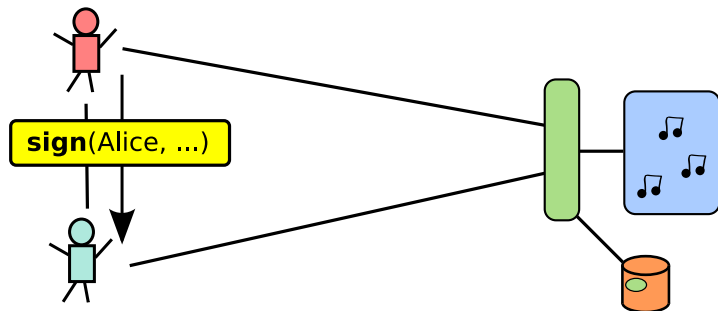
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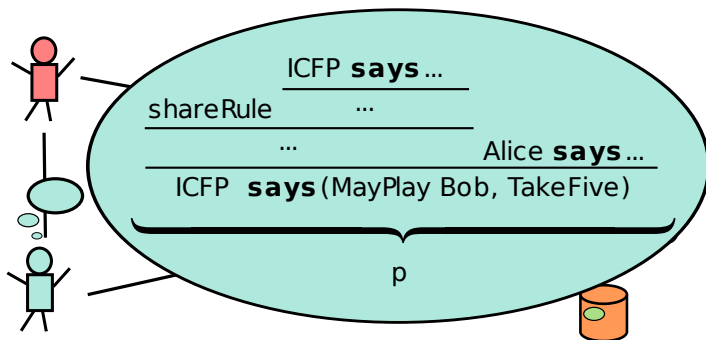
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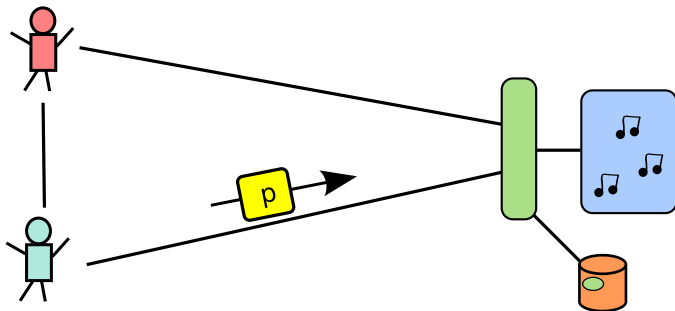
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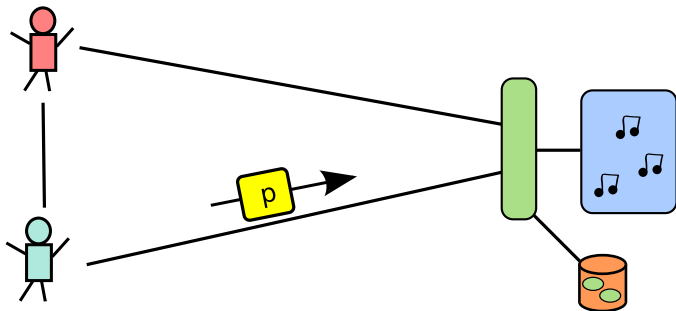
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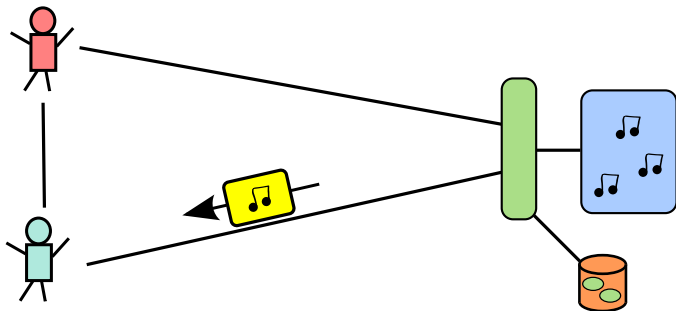
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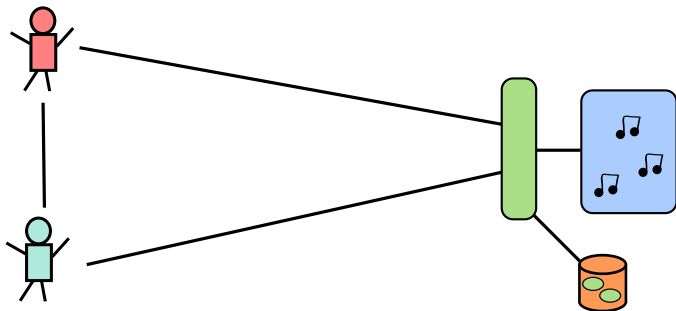
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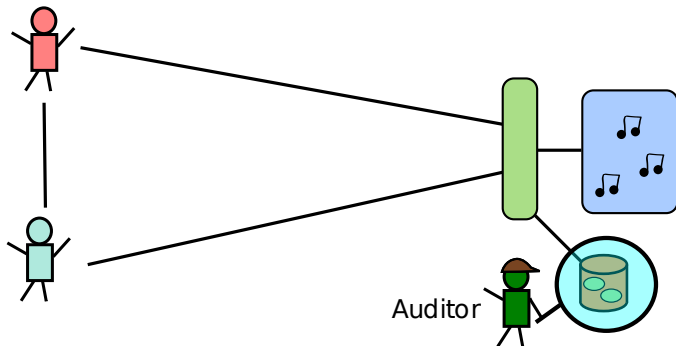
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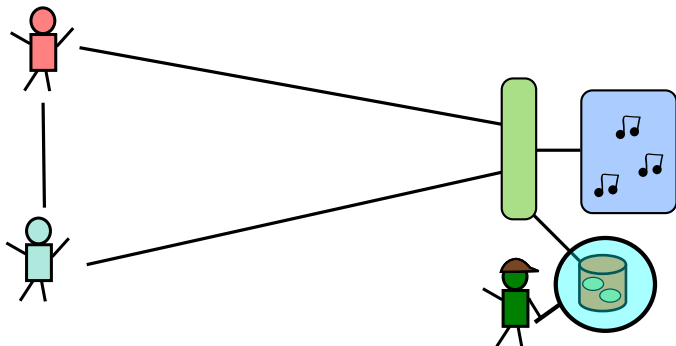
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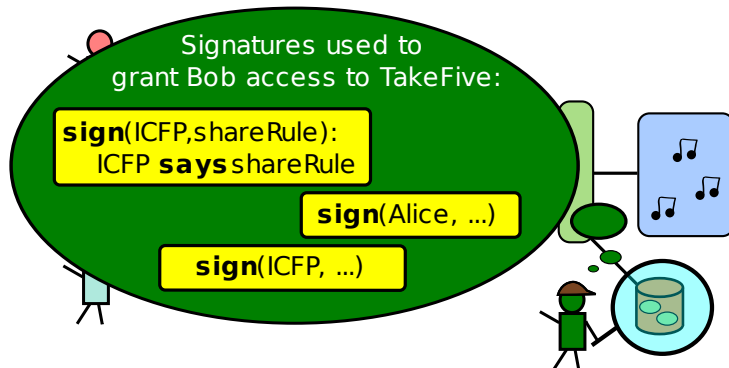
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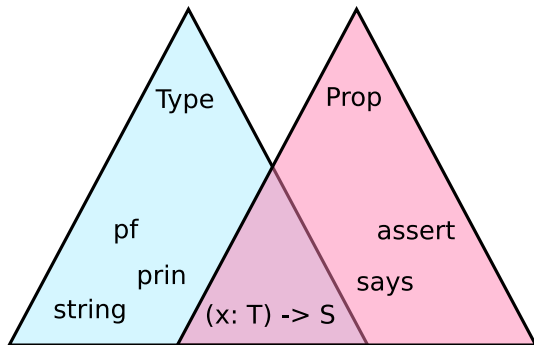
Using the ICFP policy.



Language Design and Features

AURA's type system is divided into two universes.

- Type** Contains computation expressions. Includes non-termination and world effects.
- Prop** Contains pure expressions with a clear interpretation as proofs.



Aura's says modality represents affirmation.

- The proposition “principal Alice affirms proposition P.”

Alice **says** P: **Prop**

- Principals may actively affirm propositions with signatures.

sign(Alice, P): Alice **says** P

- Principals affirm “true” propositions

return Alice p: Alice **says** P

when p: P.



DCC [Abadi+ 06], Logic with Explicit Time [DeYoung+ 08]

Dependent types allow for expressive rules.

Example (Bob acts for Alice)

Alice **says** $((P: \mathbf{Prop}) \rightarrow \text{Bob says } P \rightarrow P)$

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Bob **says** $(\text{MayPlay Bob } s) \rightarrow \text{MayPlay Bob } s)$

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Restricted formulation of dependent types:

- expressive enough for access control
- too weak for general correctness properties
- AURA feels more like ML than Coq

Effect **say** reifies a program's authority as a signature.

- Programs manufacture new **sign** objects with **say**.
- Intuitively **say** uses the program's (e.g. current user's) private key to generate the signature.
- Special principal **self** stands in for the program.

say P: **self** **says** P

say P \mapsto **sign**(**self**, P)

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say P: **pf**(**self** **says** P)

say P \mapsto **return**(**sign**(**self**, P))

Technical Point

The **pf** monad protects the **Prop** universe from **say**'s world effect.

AURA contains inductive types and *assertions*.

- Inductive Types define complex data structures.

```
data List : Type → Type {  
  | nil   : (t:Type) → (List t)  
  | cons : (t:Type) → t → (List t) → (List t) }
```

- Inductive Props define simple inference systems subject to a (draconian) positivity constraint.

```
data And: Prop → Prop → Prop {  
  | both: (P: Prop) → (Q: Prop) → P → Q → And P Q }
```

```
data False: Prop { }
```

- Assertions define access control predicates

```
assert Owns: prin → Song → Prop
```

Assertion types are uninhabited, but not false.

Inductive types admit pattern matching.

Example

```
λf: Alice says False. λP: Prop. ...  
      match f with (P) {} ...  
      : Alice says False → (P: Prop) → Alice says P
```

Assertions have no elimination form.

- Intuition: Assertions \approx type variables.
- There is no analogous function of type

ICFP **says** (Owns Bob Thriller) →
(P: **Prop**) → ICFP **says** P.

Theory and Practice

AURA's metatheory: the view from 30,000 feet.

- AURA is defined in a Pure-Type-Systems style.

$$\begin{aligned} t \quad ::= & \text{Prop} \mid \text{Type} \dots \\ & \mid (x: t) \rightarrow t \mid t \text{ says } t \dots \\ & \mid \lambda x: t. t \mid \text{sign}(t, t) \dots \end{aligned}$$

- Call-by-value reduction ensures \perp isn't confused for a proof.

Theorem (Syntactic Soundness)

Reduction preserves typing; well-typed terms don't get stuck.

Theorem (Decidability of typechecking)

Either $\Sigma; \Gamma \vdash t_1 : t_2$ or $\Sigma; \Gamma \not\vdash t_1 : t_2$, constructively.

Aura's core metatheory formalized in Coq.

- Terms *locally nameless*, with DeBruijn indexed bound variables and named free variables.
- Formalized features: inductive data types, **Prop** and **Type** language fragments, **says** and **pf** modalities. . . .

Development Size (in lines of commented Coq code)

Definitions	1400
Type Soundness	6000
Decidability of Typechecking	5000



Engineering Formal Metatheory [Aydemir+ 08]

Aura is real.

- Current Features:
 - Interpreter and typechecker for full language
 - Foreign function interface
- Coming Soon:
 - Cryptographic implementation of **sign**
 - Automatic logging
- Future Research:
 - Type inference?
 - Surface syntax?
 - Information flow?
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Demo

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Interpreter, Coq scripts, and papers available from
<http://www.cis.upenn.edu/~stevez/sol/aura.html>

- Access Control Matrices and Capabilities
- Mechanizing AURAw as a positive experience.

Conventional techniques handle the ICFP policy poorly.

Access control matrices

- ICFP server stores the list of owners and delegations.
- Owner must contact ICFP server directly to delegate.
- All participants must trust server's records re: delegation.

Atomic capabilities

- Unforgeable, atomic tokens serve as tickets to play songs.
- Who issues the tokens?

Mechanizing AURA was a positive experience.

- Aura is large.
 - 21 syntactic forms
 - 15 judgments
 - 63 inference rules
- Mechanization helped us manage AURA's complexity. Coq proofs. . .
 - provided assurance that we hadn't make mistakes.
 - enabled us to experiment without rechecking pages of unaffected proofs.
 - simplified collaboration (source control, etc.).